

# Tensor complexity and completion using hypergraph expanders

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Joint work with Yizhe Zhu, UCSD Math



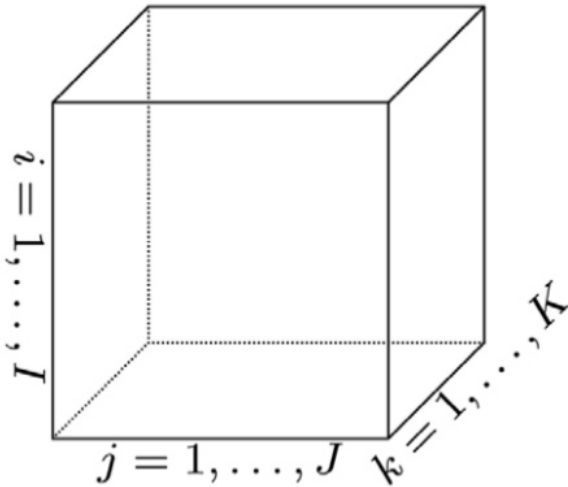
# Overview

- 1) Introduction
- 2) Tensor complexity: norms & quasinorms
- 3) Hypergraph observation model
- 4) Tensor completion bound

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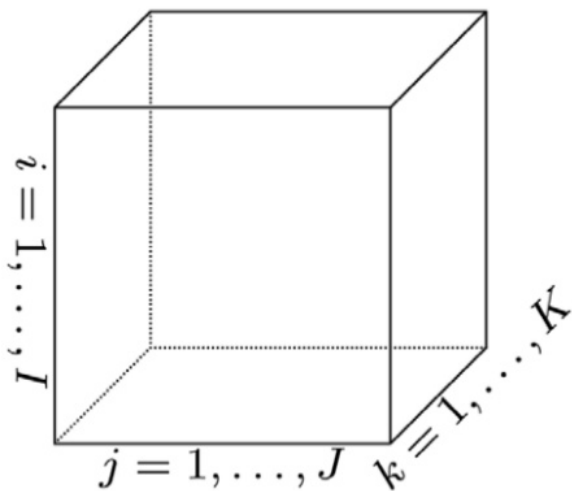
# What is a tensor?



A third-order tensor:  $\mathfrak{X} \in \mathbb{R}^{I \times J \times K}$

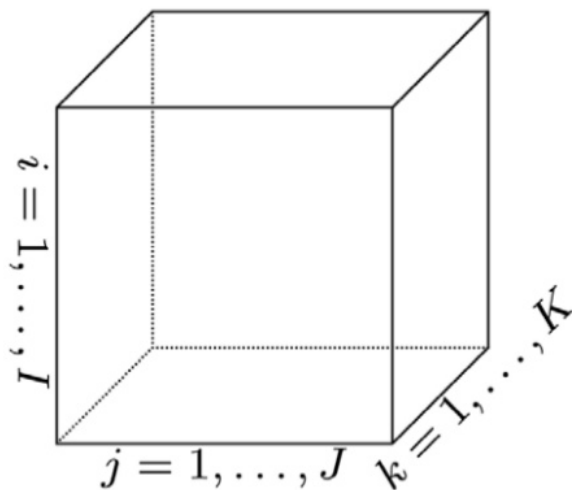
# What is a tensor?

For us, a multi-dimensional array



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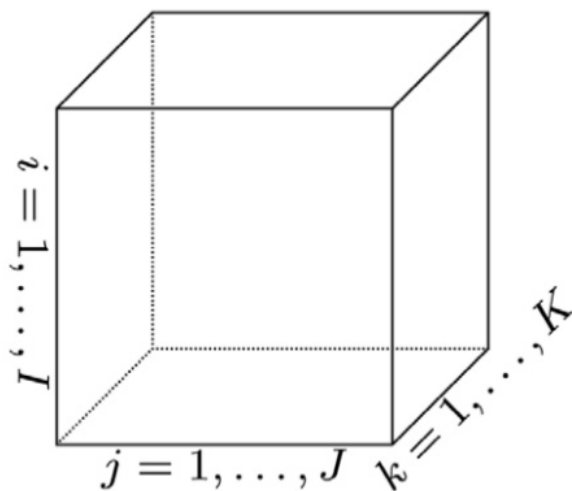


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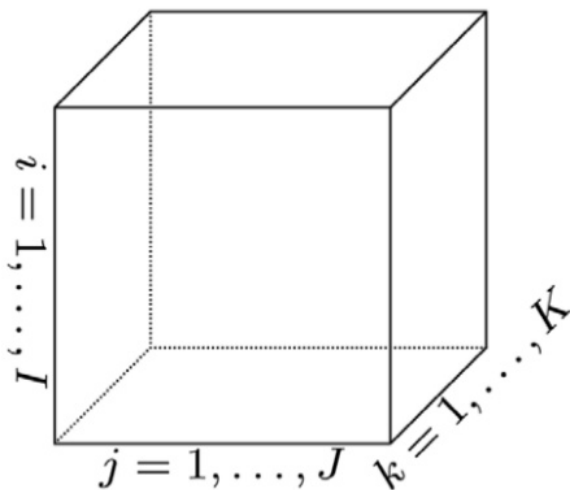
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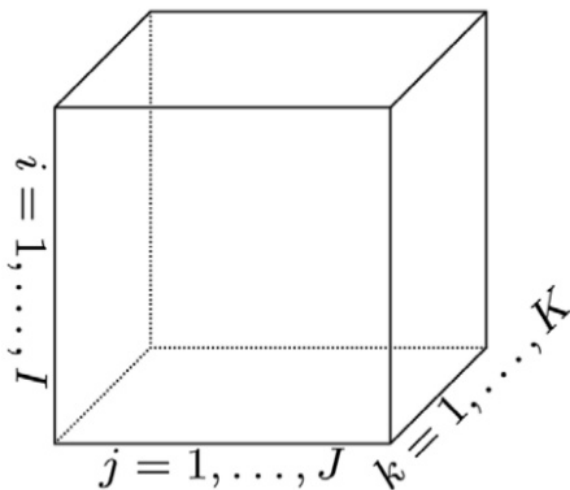
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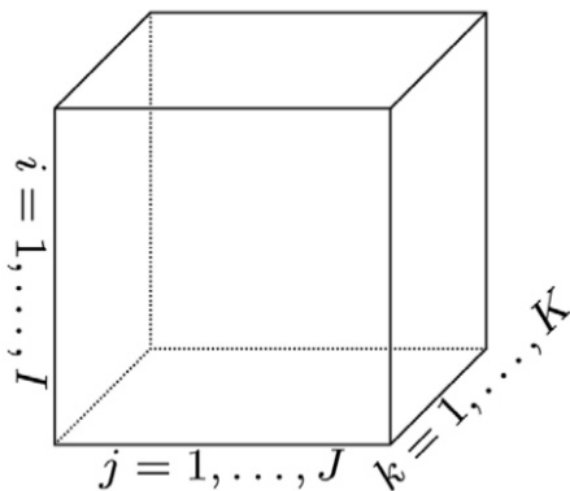
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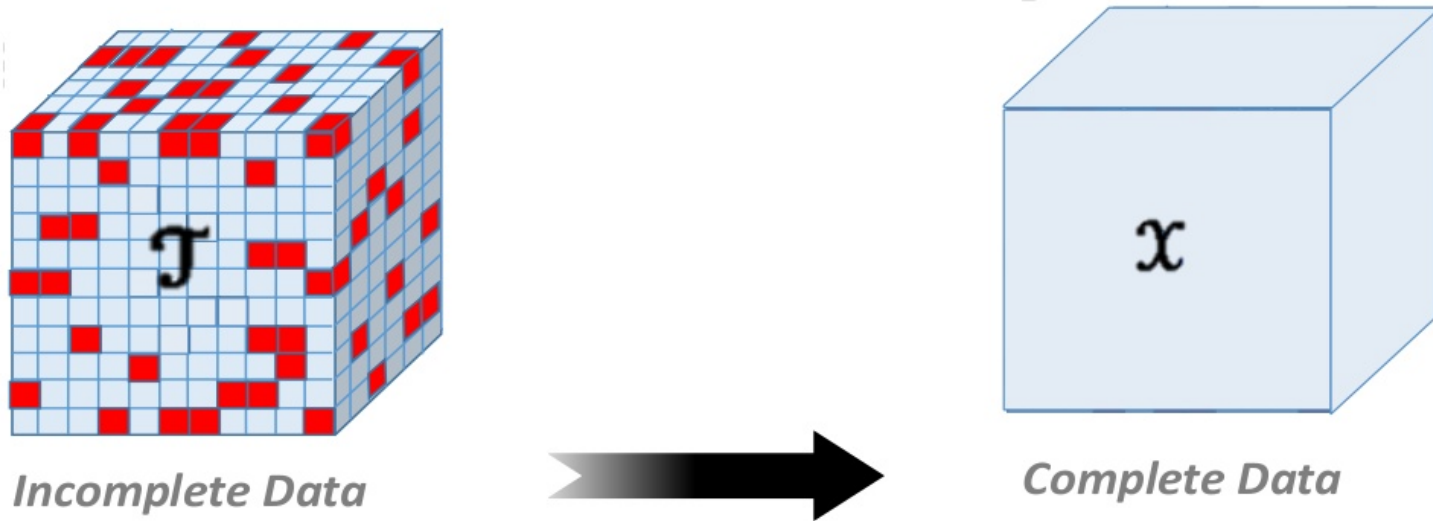
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Excellent introduction: Kolda & Bader. SIAM Rev (2009)

# Many problems cast as tensor completion

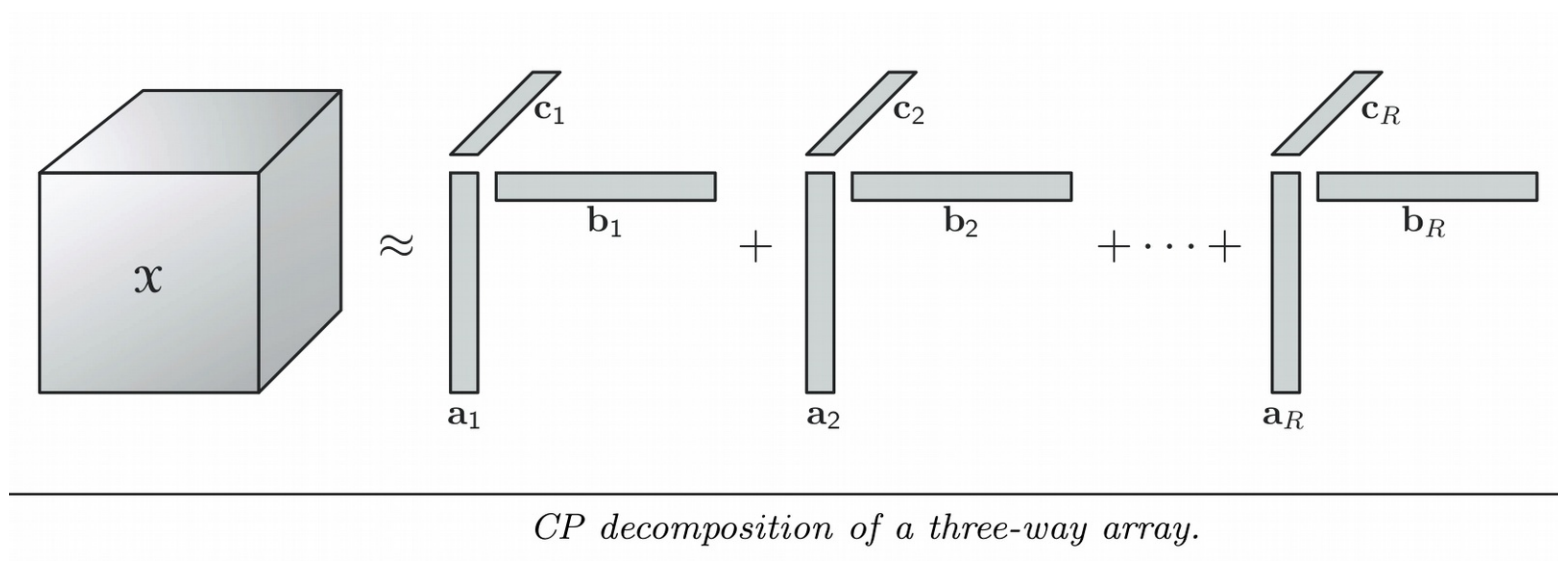
- Use low-rank structure to infer missing data



Song, Ge, Caverlee, Hu. KDD (2019)

# But tensor rank is not like in matrices

- More than one version (CP/Kruskal, Tucker, unfolding rank)
- “Most tensor problems are **NP-hard**”, Hillar & Lim. ACM (2013)



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- Matrix results:
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  - Deterministic observations? Communication complexity & rank
  - Musco, Musco, Woodruff. Arxiv (2020) ... polynomial algorithms **if** rank is relaxed, very relevant for us, but different objective

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- Deterministic bound of generalization error for completion

$$\min_T \|T\|_{\max} \text{ s.t. } \|\Omega * (T - Z)\| \leq \delta$$

# Tensor notation setup

$$T \in \bigotimes_{i=1}^t \mathbb{R}^n$$

order  $t$

$$(T)_{i_1, \dots, i_t} = \sum_{j=1}^r \underbrace{U_{i_1, j}^{(1)} \cdots U_{i_t, j}^{(t)}}_{\text{factor matrices}}$$

rank- $r$  tensor (CP),  
 $r \ n \ t = \#$  parameters

$$T = U^{(1)} \circ \dots \circ U^{(t)}$$

shorthand for rank decomposition

$$T \otimes S, \quad T * S$$

Kronecker, Hadamard products

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- Nuclear/trace norm

$$\|A\|_* = \sum_i |\sigma_i| = \min_{A=UV^\top} \|U\|_F \|V\|_F = \min_{A=UV^\top} \frac{1}{2} (\|U\|_F^2 + \|V\|_F^2)$$

Srebro & Shraibman. COLT (2005)

# Max-norm for matrices

$$\|A\|_{\max} = \min_{UV^{\mathsf{T}}=A} \|U\|_{2,\infty} \|V\|_{2,\infty} \text{ where } \|U\|_{2,\infty} = \max_i \|U_{i,:}\|_2$$

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Communication complexity & discrepancy

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Matrix completion, incoherence, leverage

**Srebro** & Shraibman 2005; Heiman et al. 2014; Cai & Zhou 2016; Foucart et al. 2017

# Max-quasinorm of a tensor

$$\|T\|_{\max} = \min_{T=U^{(1)} \circ \dots \circ U^{(t)}} \prod_{i=1}^t \|U^{(i)}\|_{2,\infty}$$

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*Lemma 1.* Let  $t \geq 2$ , then any two order- $t$  tensors  $T$  and  $S$  of the same shape satisfy the following properties:

1.  $\|T\|_{\max} = 0$  if and only if  $T = 0$ .
2.  $\|cT\|_{\max} = |c|\|T\|_{\max}$ , where  $c \in \mathcal{R}$ .
3.  $\|T + S\|_{\max} \leq \underbrace{\left(\|T\|_{\max}^{2/t} + \|S\|_{\max}^{2/t}\right)^{t/2}}_{\substack{p\text{-norm, } p = 2/t \\ \text{Dilworth (1985)}}} \underbrace{\leq 2^{t/2-1} (\|T\|_{\max} + \|S\|_{\max})}_{\text{quasi-triangle inequality}}.$

# New results for max-qnorm

**Theorem 1.** *Let  $T \in \bigotimes_{i=1}^t n_i$  and  $S \in \bigotimes_{i=1}^t m_i$ , then:*

1.  $\|T_{I_1, \dots, I_t}\|_{\max} \leq \|T\|_{\max}$  for any subsets of indices  $I_i \subseteq [n_i]$

2.  $\|T \otimes S\|_{\max} \leq \|T\|_{\max} \|S\|_{\max}$

3.  $\|T * S\|_{\max} \leq \|T \otimes S\|_{\max}$ , where  $T, S \in \bigotimes_{i=1}^t n_i$

4.  $\|T * T\|_{\max} \leq \|T\|_{\max}^2$

Generalizes a number of results in  
“A Direct Product Theorem for Discrepancy”. Lee, Shraibman, Špalek (2008)

# Comparison to $\ell_\infty^{n^t}$

Lower bound:  $\|T\|_{\max} \geq \max_{1 \leq i \leq t} \|T_{[i]}\|_{\max} \geq |T|_\infty$

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**A guess:**  $\|T\|_{\max} \leq \sqrt{r^{t-1}} \cdot |T|_\infty$

# Sign tensors & another norm

$$\begin{bmatrix} + & - & + \\ - & + & - \\ + & - & + \end{bmatrix} = \begin{bmatrix} + \\ - \\ + \end{bmatrix} \begin{bmatrix} + & - & + \end{bmatrix} \quad \text{rank-1 sign tensor}$$

**Sign nuclear norm:**

$$\|T\|_{\pm} = \inf \left\{ \sum_{i=1}^r |\alpha_i| \mid T = \sum_{i=1}^r \alpha_i S_i \text{ where } \alpha_i \in \mathbb{R}, \text{rank}_{\pm}(S_i) = 1 \right\}$$

# Relation between sign and max-qnorm

**Lemma 2.** *The sign nuclear norm and max-qnorm satisfy*

$$\|T\|_{\pm} \leq K_G^{t-1} \|T\|_{\max},$$

*where  $K_G$  is the Grothendieck constant over  $\mathbb{R}$ .*

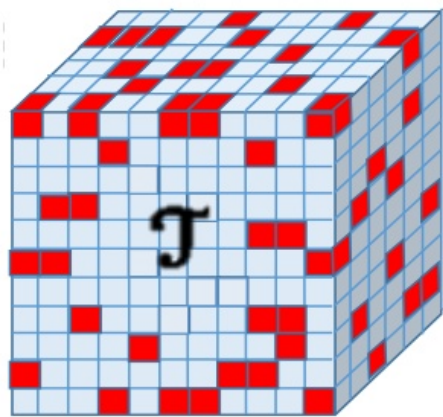
(Tightens a result by Ghadermarzy et al.)



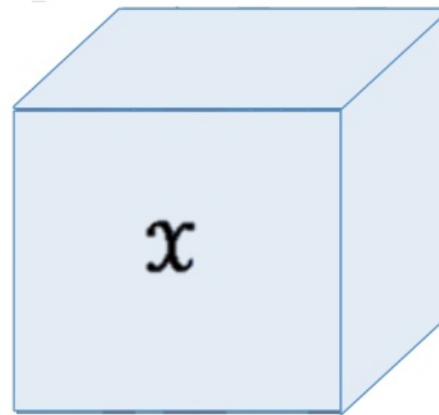
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# Modeling observations



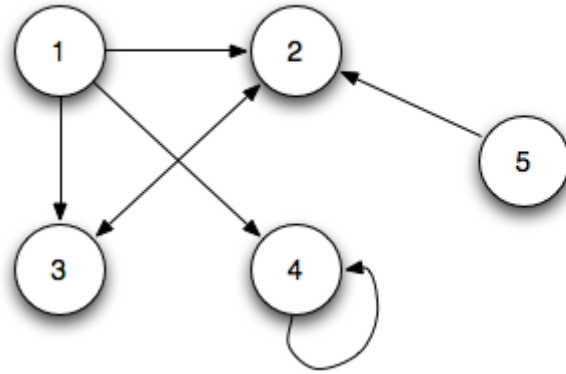
*Incomplete Data*



*Complete Data*

sparse binary mask

# Adjacency matrix of a graph

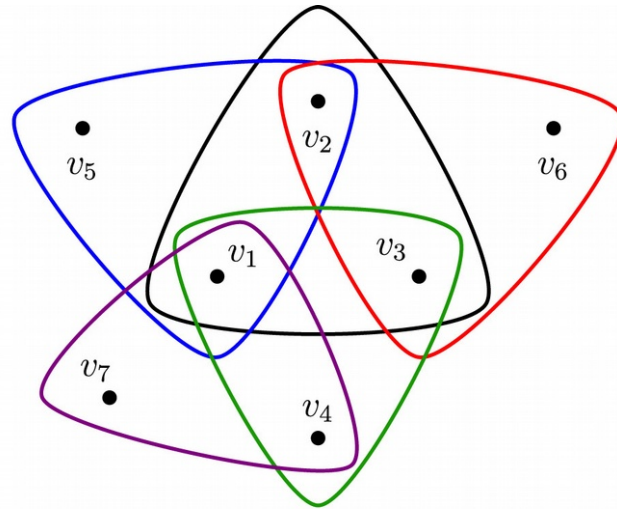


	1	2	3	4	5
1	0	1	1	1	0
2	0	0	1	0	0
3	0	1	0	0	0
4	0	0	0	1	0
5	0	1	0	0	0

# Hypergraph $\rightarrow$ Adjacency tensor

Properties we require

- $t$ -uniform: all hyperedges contain  $t$  vertices
- $t$ -partite: non-symmetric tensors



Lofty goal:  $H$  has “good” mixing

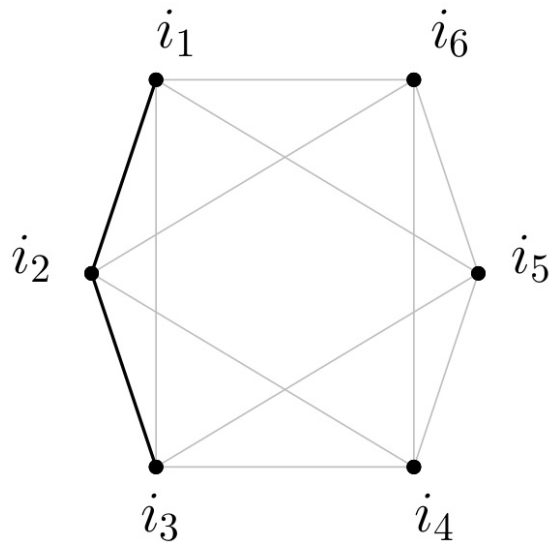
$$| e(V_1, \dots, V_t) - \alpha |V_1| \cdots |V_t| | \leq \lambda \sqrt{|V_1| \cdots |V_t|}$$

$$\lambda = \|T - \alpha J\|_\sigma$$

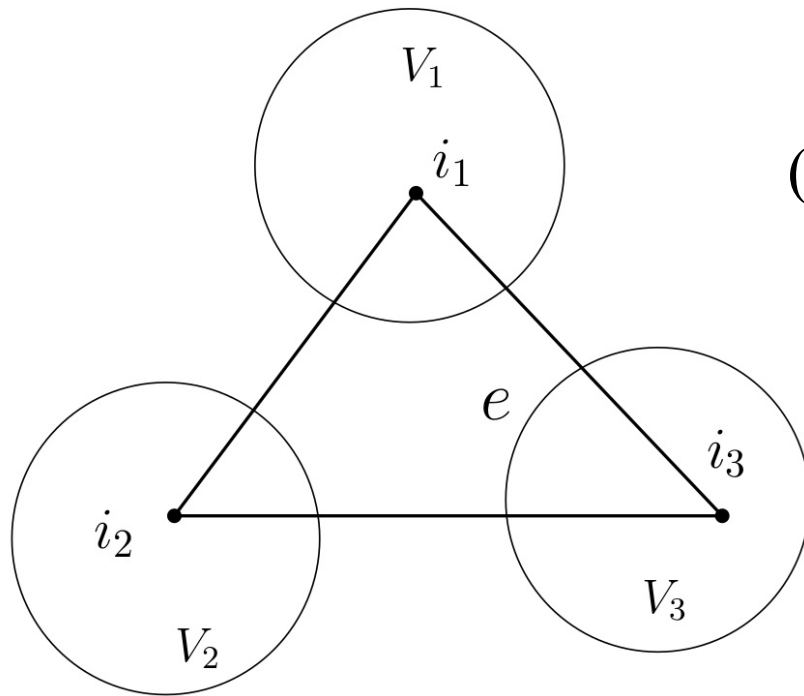
“Second eigenvalue” of  $T(H)$

$$\|T\|_\sigma = \sup_{v_1, \dots, v_t \in S^{n-1}} \left| \sum_{i_1, \dots, i_t=1}^n T_{i_1, \dots, i_t} v_1(i_1) \cdots v_t(i_t) \right|$$

# Expander construction



$G$



$(d = 4, t = 3)$

$H$

Alon et al. Computational Complexity (1995)  
Bilu & Hoory. EJ Combinatorics (2004)

has  $nd^{t-1}$  many edges

# (Weak) mixing lemma

**Lemma 3.** *Construct  $H$  from the  $d$ -regular base graph  $G$ . Then:*

$$\left| \frac{e(W_1, \dots, W_t)}{nd^{t-1}} - \prod_{i=1}^t \alpha_i \right| \leq \left( \left( 1 + \frac{\lambda}{d} \right)^{t-1} - 1 \right) \min \left\{ \frac{1}{4}, \prod_{i=1}^t \max \{ \sqrt{\alpha_i}, \sqrt{1 - \alpha_i} \} \right\}$$

where

$$\alpha_i = \frac{|W_i|}{n} \quad \text{and} \quad \lambda = \lambda_2(G).$$

Tightens results from Alon et al. (1995); Bilu & Hoory (2004)

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# Max-qnorm penalized guarantee

- Suppose we can *solve*

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## Theorem 2:

$$\frac{1}{n^t} \|\hat{T} - T\|_F^2 \leq (4K_G)^{t-1} \|T\|_{\max}^2 \left( \left(1 + \frac{\lambda}{d}\right)^{t-1} - 1 \right) + 4\delta^2$$

Ghademarzy, Plan, Yilmaz (2018)

Heiman et al. (2014)

Brito, Dumitriu, Harris (2018)

# Sample complexity

- # observations  $|E(H)| = nd^{t-1}$
- use an expander  $G \implies \frac{\lambda}{d} = \frac{1}{\sqrt{d}}$

Required # samples:

$$|E| = n \left( \frac{(t-1)(4K_G)^{t-1} \|T\|_{\max}^2}{\varepsilon} \right)^{2(t-1)}$$

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- 3) Use mixing lemma to bound variation from sample

$$\left| \frac{1}{n^t} \sum_{e \in [n]^t} Q_e - \frac{1}{nd^{t-1}} \sum_{e \in E} Q_e \right| \leq 2^{t-2} \left( \left( 1 + \frac{\lambda}{d} \right)^{t-1} - 1 \right) \|Q\|_{\pm}$$

# Proof sketch (part 2)

1) Take  $Q = (\hat{T} - T) * (\hat{T} - T)$  squared residuals

2) Then we have:

$$\begin{aligned}\|Q\|_{\pm} &\leq 2^{t-2} K_G^{t-1} \left( \|\hat{T}\|_{\max} + \|T\|_{\max} \right)^2 \\ &\leq 2^t K_G^{t-1} \|T\|_{\max}^2\end{aligned}$$

# Conclusions

- New inequalities for tensor complexity measures
- Better weak mixing for expander hypergraph
- New analysis of tensor completion
  - Linear dependence on  $n$
  - Better bound possible with 2<sup>nd</sup> eigenvalue of  $H$

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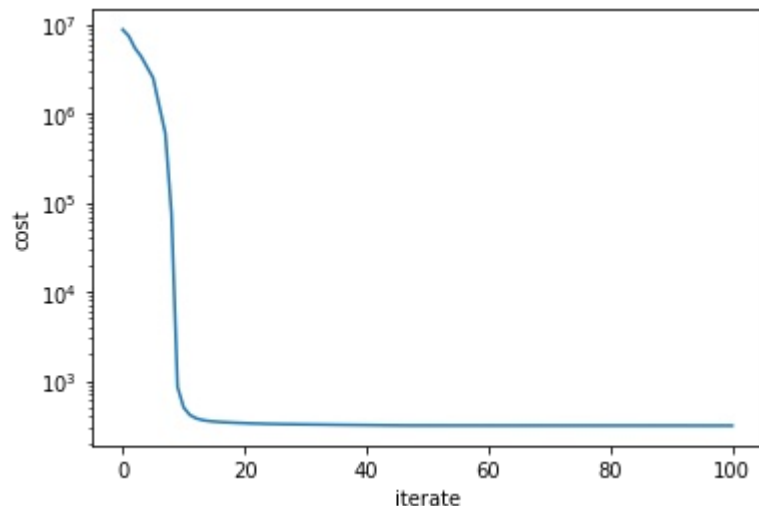
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- Approx solution, get bound w/  $\max(\|\hat{T}\|_{\max}, \|T\|_{\max})$
- Suggests even approx solutions hard!

# Early numerical experiments



Coordinate descent  
on factors  $U^{(i)}$

- $r > \text{rank}(T)$  helps
- overfit residuals okay
- stagnation for low rank,  
too few observations

$$n = 200, t = 3, r = 4, |E| = 55224$$
$$\sim 0.7\%$$

# Thank you!

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  - **coauthor: Yizhe Zhu**
  - Ioana Dumitriu
  - Paul Beame
  - Adi Shraibman
- Funding:
  - NSF DMS-1712630 (YZ)



“Deterministic tensor completion  
with hypergraph expanders”

<https://arxiv.org/abs/1910.10692>

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- Let  $Z = \text{span} \left\{ U_{:,i}^{(1)} \circ \cdots \circ U_{:,i}^{(t)} \right\}_{i=1}^r$ ,  $X = \left( \mathbb{R}^{n^t}, \|\cdot\|_{\max} \right) \cap Z$ ,  $Y = \left( \mathbb{R}^{n^t}, |\cdot|_{\infty} \right) \cap Z$

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- By Dilworth (1985),  $r$ -dim  $p$ -norm space  $d(X, \ell_2) \leq r^{\frac{2}{p} - \frac{1}{2}}$
- Let  $Z = \text{span} \left\{ U_{:,i}^{(1)} \circ \dots \circ U_{:,i}^{(t)} \right\}_{i=1}^r$ ,  $X = (\mathbb{R}^{n^t}, \|\cdot\|_{\max}) \cap Z$ ,  $Y = (\mathbb{R}^{n^t}, |\cdot|_{\infty}) \cap Z$
- Then  $d(X, Y) \leq d(X, \ell_2) d(Y, \ell_2) \leq \sqrt{r^t}$



# Ideas for $\|T\|_{\max} \leq \sqrt{r^t} \cdot |T|_{\infty}$

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- Consider norm inequality  $c|T|_{\infty} \leq \|T\|_{\max} \leq C|T|_{\infty}$